

Logic and Artificial Intelligence

Lecture 17

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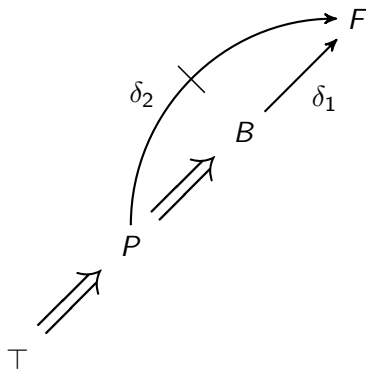
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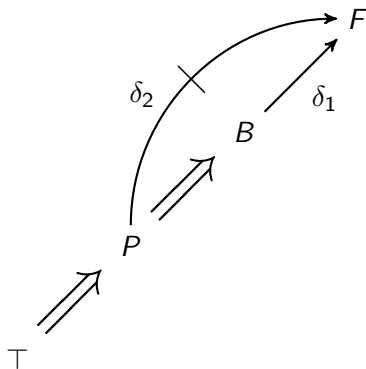
A Distinction

Rebutting vs. undercutting

Tweety Triangle: Rebutting



Tweety Triangle: Rebutting



δ_1 is the default: $B \longrightarrow F$

δ_2 is the default $B \longrightarrow \neg F$

There is a *conflict* in that $Conclusion(\delta_1)$ is inconsistent with $Conclusion(\delta_2)$

Undercutting, I

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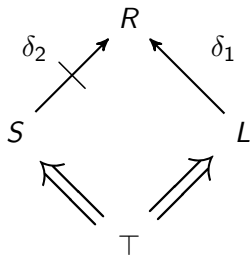
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$$S \longrightarrow \neg R$$



$$\delta_1 < \delta_2$$

$$\delta_2 < \delta_1$$

Undercutting, II

If the object looks red but the reliable source tells me otherwise, then it is natural to appeal to another default, with the conclusion that the object is not red, since what the reliable source says tends to be true and the reliable source has told me that it is not red. And because, by hypothesis, the reliable source is more reliable than perception, this new default would have to be stronger than the original, that whatever looks red tends to be red, and so would defeat this original default in the sense we have considered so far, by providing a stronger reason for a conflicting conclusion.

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$\delta_2 < \delta_1$, so the agent believes the object is not red.

Undercutting, III

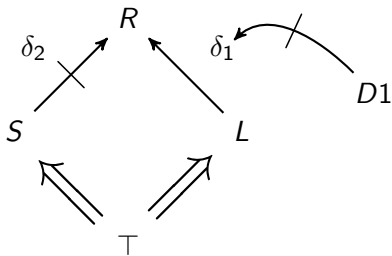
Suppose that I have taken a drug — let us call it Drug 1 — that makes everything look red. If the object looks red but I have taken Drug 1, on the other hand, then it seems again that I am no longer entitled to the conclusion that the object is red. But in this case, the original default is not defeated in the same way. There is no stronger reason for concluding that the object is not red; instead, it is as if the favoring relation represented by the original default is itself severed, so that what was once a reason no longer provides any support for its conclusion.

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$$D1 \longrightarrow \text{Out}(\delta_1)$$

Undercutting, IV



When should a rational agent believe/accept/assert and *conditional*? (If A , then B : $A > B$)

Note: $A > B$ is NOT the same as $A \Rightarrow B$

Ramsey Test

If two people are arguing 'If p , then q ?' and are both in doubt as to p , they are adding p hypothetically to their stock of knowledge and arguing on that basis about q ; so that in a sense 'If p , q ' and 'If p , $\neg q$ ' are contradictories. We can say that they are fixing their degree of belief in q given p . If p turns out false, these degrees of belief are rendered void. If either party believes not p for certain, the question ceases to mean anything to him except as a question about what follows from certain laws or hypotheses.

F. P. Ramsey. *General Propositions and Causality*. 1929.

AGM and the Ramsey Test

Recall that \mathcal{L} is a propositional language. Let $\mathcal{L} \subseteq \mathcal{L}_<$ extends \mathcal{L} with a conditional operator ' $>$ ': $A > B$ means 'If A , B '

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Belief revision model: $\langle \mathbf{K}, * \rangle$ where $* : \mathbf{K} \times \mathcal{L} \rightarrow \mathbf{K}$.

Gärdenfors Impossibility Result, I

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Gärdenfors Impossibility Result, II

From (R) we can conclude

(M) For all belief sets K , K' and propositions A , $K \subseteq K'$ implies $K_A \subseteq K'_A$

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Preservation Criterion

(P) If $\neg A \notin K$ and $B \in K$ then $B \in K_A$

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From (1) and (P)

$$(I) \text{ If } \neg A \notin K \text{ then } K/A \subseteq K_A$$

Gärdenfors Impossibility Result, IV

Two propositions A and B are *disjoint* provided $\vdash \neg(A \wedge B)$

A belief revision model $\langle \mathbf{K}, \mathbf{F} \rangle$ is **nontrivial** provided there is a $K \in \mathbf{K}$ and *pairwise disjoint* propositions A, B and C such that $\neg A \notin K$, $\neg B \notin K$ and $\neg C \notin K$.

Theorem There is no nontrivial belief revision model that satisfies all the conditions (1), (2), (M) and (P).

P. Gärdenfors. *Belief Revision and the Ramsey Test*. *Philosophical Review*, Vol. 95, pp. 81 - 93, 1986.

Gärdenfors Impossibility Result: Proof

Nontrivial There are A , B , and C such that $\vdash \neg(A \wedge B)$,
 $\vdash \neg(B \wedge C)$, $\vdash \neg(A \wedge C)$ and belief set K such that $\neg A \notin K$,
 $\neg B \notin K$ and $\neg C \notin K$.

(1) $A \in K_A$

(2) If $K \neq K_\perp$ and $K_A = K_\perp$ then $\vdash \neg A$

(R) $A > B \in K$ iff $B \in K * A$

(M) For all K , K' and A , $K \subseteq K'$ implies $K_A \subseteq K'_A$

(P) If $\neg A \notin K$ and $B \in K$ then $B \in K_A$

(3) $(K/A)/B = K/(A \wedge B)$

(4) $K/(A \vee B) \subseteq K/A$

(I) If $\neg A \notin K$ then $K/A \subseteq K_A$

H. Leitgeb. *On the Ramsey Test without Triviality*. Notre Dame Journal of Formal Logic, 51:1, 2010, pp. 21 - 54.

H. Leitgeb. *Beliefs in Conditionals Vs. Conditional Beliefs*. Topoi, 26 (1), 2007, pp. 115 - 132.

I. Levi. *Iteration of Conditionals and the Ramsey Test*. Synthese, 76, 1988, pp. 49 - 81.

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- ▶ Static vs. dynamic

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- ⇒ informational attitudes (eg., knowledge, belief, certainty)
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