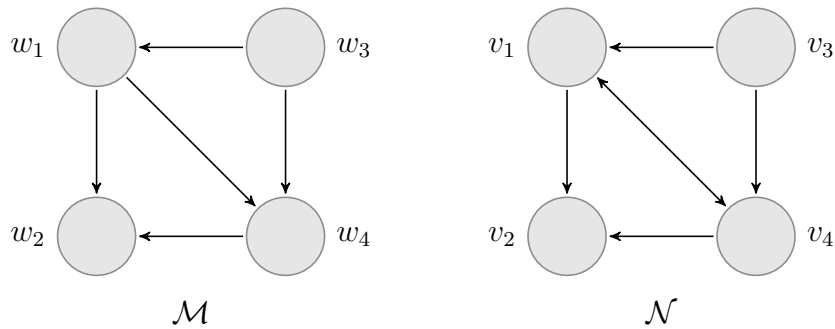


## Homework #3

(all questions have equal weight)

1. (Compare this exercise with exercise 6(a) in Section 1.2 of Enderton). Let  $\mathcal{M} = \langle W, R, V \rangle$  and  $\mathcal{M}' = \langle W, R, V' \rangle$  be two models that differ only in their valuation functions. Suppose that  $\varphi$  is a modal formula where  $V(w, p) = V'(w, p)$  for all states and sentence letters  $p$  occurring in  $\varphi$ . Prove that  $\mathcal{M} \models \varphi$  iff  $\mathcal{M}' \models \varphi$ . [Hint: the proof is by induction on the structure of  $\varphi$ .]
  
2. Consider the following two relational structures:

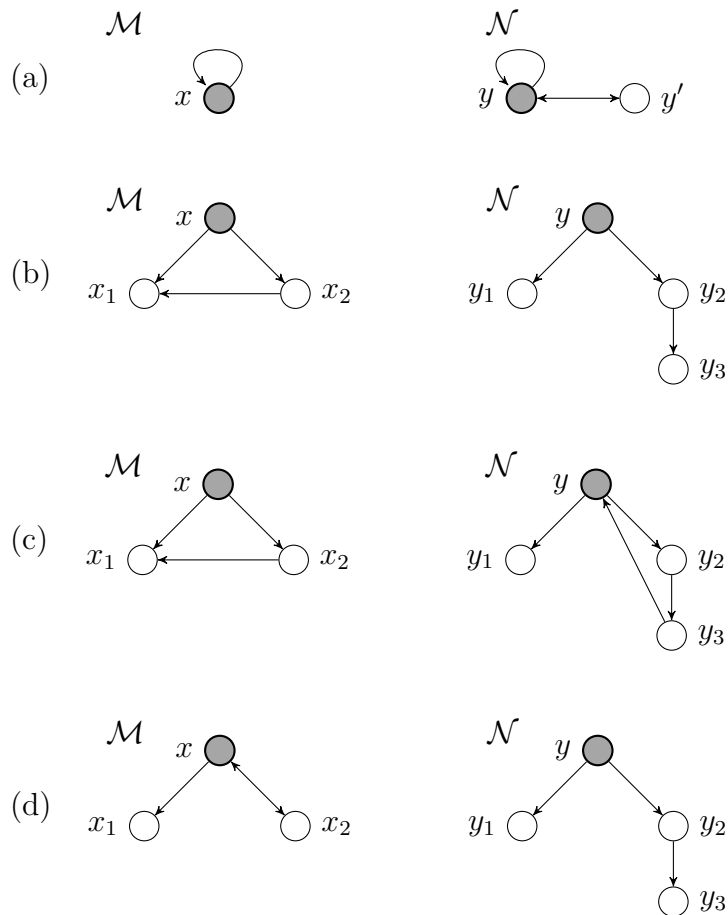


- (a) Show that the following sets are definable in  $\mathcal{M}$ :  $\emptyset$ ,  $\{w_1\}$ ,  $\{w_2\}$ ,  $\{w_3\}$ ,  $\{w_4\}$ ,  $\{w_1, w_2, w_3, w_4\}$ .
- (b) Which of the following sets are definable in  $\mathcal{N}$ :  $\{v_1\}$ ,  $\{v_1, v_3\}$ ,  $\{v_1, v_2, v_4\}$ ? (Note: you must explain your answers.)
  
3. (a) Prove that  $\diamond(P \rightarrow Q) \rightarrow (\Box P \rightarrow \diamond Q)$  is valid.  
 (b) Show that  $\Box(\Box P \rightarrow Q) \vee \Box(\Box Q \rightarrow P)$  is not valid.
  
4. Consider the following new operator  $\Box^\leftarrow$  defined on relational structures  $\mathcal{M} = \langle W, R, V \rangle$  as follows:

$$\mathcal{M}, w \models \Box^\leftarrow \varphi \text{ iff for all } v \in W, \text{ if } vRw \text{ then } \mathcal{M}, v \models \varphi$$

Prove that  $\Box^\leftarrow$  is not definable in the modal language.

5. Is there a bisimulation connecting the states  $x$  and  $y$  in the following models? Justify your answer (that is, if the answer is yes, then provide a bisimulation; and if the answer is no, then find a modal formula that distinguishes the two states).



6. A frame is a pair  $\langle W, R \rangle$  where  $W$  is a nonempty set and  $R \subseteq W \times W$  is a relation. We say that a formula is **valid on a frame**  $\mathcal{F} = \langle W, R \rangle$ , denote  $\mathcal{F} \models \varphi$ , provided  $\mathcal{F}, V, w \models \varphi$  for all states  $w \in W$  and valuations  $V$ .

- (a) We say  $\mathcal{F} = \langle W, R \rangle$  is **Euclidean** if  $R$  is a Euclidean relation (for all  $x, y, z \in W$  if  $xRy$  and  $xRz$  then  $yRz$ ). Prove that  $\mathcal{F}$  is Euclidean iff  $\mathcal{F} \models \diamond\varphi \rightarrow \square\diamond\varphi$ .

- (b) We say  $\mathcal{F} = \langle W, R \rangle$  is **dense** if  $R$  is a dense relation (for all  $x, y, z \in W$  if  $xRy$ , then there is a  $z \in W$  such that  $xRz$  and  $zRy$ ). Prove that  $\mathcal{F}$  is dense iff  $\mathcal{F} \models \Box\Box\varphi \rightarrow \Box\varphi$ .

**The homework is DUE Friday, January 30 at 10AM in class.**